

UBC Math Circle 2018 Problem Set 6

Big Hint: The theme this week is invariants, monotonicity, and colouring.

1. In a group of n people, each person has at most 3 enemies (if A is an enemy of B , then B is also an enemy of A). Prove that we can split the n people into two groups such that each person has at most one enemy in her group.

Hint: Is there a way to quantify the goodness of a grouping? How can we improve the quality of a bad grouping by moving people around?

Solution: Let S be the total of enemies in the same group. Observe that whenever we move a person from a group with two or more enemies to the other group (with at most one enemy), S strictly decreases. Also observe that S is bounded below, so we can only make a finite number of these swaps before each person only has one enemy in her group.

2. Let $x_1, x_2, \dots, x_{2018}$ be integers, and $f_0(i) = x_i$. For $n > 0$, define $f_n(i)$ to be the number of integers j where $f_{n-1}(j) = f_{n-1}(i)$ (including $j = i$). Prove that there exists some integer m such that $f_m(i) = f_{m-1}(i)$ for all integers i between 1 and 2018 inclusive.

Hint: What can we say about the number of distinct values for each f_n .

Solution: Solution Sketch:

Let $g(n)$ be the number of different values of $f_n(i)$ over all i . Observe that $g(n)$ is bounded below by 1, and is non-increasing. Therefore, there exists some integer k , and some integer N , such that $g(n) = k$ for all $n \geq N$.

Suppose that $g(n-1) = g(n) = g(n+1)$, then we can show that $f_n(i) = f_{n+1}(i)$.

3. Two thousand 1×1 cells of a 2018×2018 square are infected. Each second, the cells with at least two infected neighbours become infected. Two cells are neighbours if they have a common side. Can the infection spread to the whole square?

Solution: We observe that the perimeter of the total infected area is non-increasing. We can see this by looking at what happens when a cell becomes infected. Initially the perimeter is 4×2000 , but the square has perimeter 4×2018 , so the infection will never spread to the whole square.

4. We start with the number 7^{2018} (in base 10). We repeatedly remove the first digit and add it to the remaining number until we obtain a number with 10 digits. Prove that this number has two equal digits.

Hint: What is the invariant?

Solution: It's a good guess that the digital sum is an invariant.

Observe that $7^{2018} \not\equiv 0 \pmod{9}$. We can write

$$7^{2018} = d_0 + 10^1 d_1 + 10^2 d_2 + \cdots + 10^m d_m \equiv d_0 + d_1 + \cdots + d_m \pmod{9},$$

where d_0, d_1, \dots are the digits of 7^{2018} . Observe that this means the digital sum is invariant, because

$$d_0 + 10^1 d_1 + \cdots + 10^m d_m \equiv d_0 + 10^1 d_1 + \cdots + 10^{m-1} d_{m-1} + d_m \pmod{9}.$$

Now, the sum of the ten digits in the end must also be congruent to 7^{2018} modulo 9. However, if they are all distinct, the sum will be $0 + 1 + 2 + \cdots + 9 = 45 \equiv 0 \pmod{9}$, which is not equivalent to 7^{2018} .

5. A 23×23 square grid is completely tiled with 1×1 , 2×2 , and 3×3 tiles. What is the minimum number of 1×1 tiles needed?

Hint: How can we colour the grid so that we can say stuff about the 2×2 and 3×3 tiles?

Solution: Observe that we could tile the square grid with at most one 1×1 tile. Put the 1×1 tile in the center, then tile each of the four 11×12 rectangles with a row of six 2×2 tiles, and three rows of 3×3 tiles.

Now we show that at least one 1×1 tile is needed. Suppose that no 1×1 tiles are needed. Colour the rows of the grid red and blue alternating, starting with red. Then we have 23 more red squares than blue ones. Each 2×2 tile covers the same number of red and blue squares. Each 3×3 tile covers three more of one colour. This means the difference between the red and blue squares should be a multiple of 3. 23 is not a multiple of three, so we can't tile the grid with only 2×2 and 3×3 tiles.

6. Rachel has n one's are written on the blackboard. Each step, Rachel erases two numbers a and b and replaces them with $\frac{a+b}{4}$. She repeats this until there is only one number left on the blackboard. Prove that this number is at least $\frac{1}{n}$.

Solution: The *AM-GM-HM* inequality gives

$$\frac{a+b}{4} \geq \frac{1}{2} \cdot \frac{2}{\frac{1}{a} + \frac{1}{b}} \implies \frac{1}{\left(\frac{a+b}{4}\right)} \leq \frac{1}{a} + \frac{1}{b}.$$

This means that the sum of the reciprocals of the numbers on the blackboard is non-increasing.

Let the final number be x . Initially the sum of reciprocals is n , so $\frac{1}{x} \leq n$. Therefore $x \geq \frac{1}{n}$.

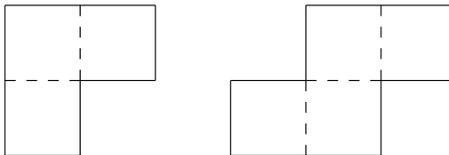
7. Eric placed m cookies at the vertices of a regular n -gon, where $m > n$. Each step, Eric chooses a vertex with at least two cookies, moves one cookie to the clockwise neighbour, and another to the counterclockwise neighbour. After k moves, the original arrangement of cookies is restored. Prove that k is a multiple of n .

Solution: Label the vertices in order from 1 to n . Let a_i be the number of times Eric chose vertex i . Observe that each time Eric chose vertex i , two cookies were removed from vertex i , and each time Eric chose one of the two vertices adjacent to vertex i , one cookie is added to vertex i . Since the original arrangement is restored, the number of cookies removed and added to each vertex must be equal. Therefore, we get the following equations

$$\begin{aligned} 2a_1 &= a_n + a_2 \\ 2a_2 &= a_1 + a_3 \\ &\vdots \\ 2a_n &= a_{n-1} + a_1 \end{aligned}$$

Without loss of generality, let a_1 be the maximum of all a_i . Then $a_2 = a_n = a_1$ from the first equation. It follows that $a_1 = a_2 = \dots = a_n$, so the number of moves is equal to na_1 , which is always a multiple of n .

8. (Putnam 2016 A4). Consider a $(2m - 1) \times (2n - 1)$ rectangular region, where m and n are integers such that $m, n \geq 4$. This region is to be tiled using tiles of the two types shown:



What is the minimum number of tiles required to tile the region?

Solution: Putnam 2016 A4.

Hint: We colour the grid red and blue such that each tile covers at most one red square. This is a lower bound on the number of tiles. We can then show that this lower bound is achievable by construction.

<http://kskedlaya.org/putnam-archive/>